

# Kleene logic and inference

Grzegorz Malinowski

Department of Logic  
Łódź University

2013

**The aims:**

## **The aims:**

- analysis of a distinguished three-valued 1938 Kleene logic construction

## The aims:

- analysis of a distinguished three-valued 1938 Kleene logic construction
- outlining the problem of inference, as important in view of the fact that no formula of Kleene logic is a  $K_3$  tautology

## The aims:

- analysis of a distinguished three-valued 1938 Kleene logic construction
- outlining the problem of inference, as important in view of the fact that no formula of Kleene logic is a  $K_3$  tautology
- critical discussion of 1974 Cleve's attempt to define a truth-preserving consequence relation within a model theoretic framework for inexact predicates

## The aims:

- analysis of a distinguished three-valued 1938 Kleene logic construction
- outlining the problem of inference, as important in view of the fact that no formula of Kleene logic is a  $K_3$  tautology
- critical discussion of 1974 Cleve's attempt to define a truth-preserving consequence relation within a model theoretic framework for inexact predicates
- showing a radically different and sound ground for Kleene inference supported by two "idealization" operators which convert neutral propositions into false or true

## 1. Kleene logic

## 1. Kleene logic

Kleene three-valued logic [1938], [1952] has an epistemological motivation. The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation.

## 1. Kleene logic

Kleene three-valued logic [1938], [1952] has an epistemological motivation. The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation.

- is inspired by the studies of the foundations of mathematics, especially by the problems of algorithms, cf. also Kleene [1952]

## 1. Kleene logic

Kleene three-valued logic [1938], [1952] has an epistemological motivation. The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation.

- is inspired by the studies of the foundations of mathematics, especially by the problems of algorithms, cf. also Kleene [1952]
- is based on the third category of propositions, whose logical value of truth or falsity is not essential, undefined or undetermined by means of accessible algorithms

The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation.  
And,

The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation.

And,

- its assumed status is somewhat different from the classical values of truth, t and falsity, f

The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation.  
And,

- its assumed status is somewhat different from the classical values of truth, t and falsity, f
- it is not independent since it may turn into either of the two classical values

The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation. And,

- its assumed status is somewhat different from the classical values of truth, t and falsity, f
- it is not independent since it may turn into either of the two classical values
- the set f, u, t of three values cannot be considered as linearly ordered

The third logical value was designed to mark indeterminacy of some proposition at a certain stage of scientific investigation. And,

- its assumed status is somewhat different from the classical values of truth, t and falsity, f
- it is not independent since it may turn into either of the two classical values
- the set f, u, t of three values cannot be considered as linearly ordered

The connectives:  $\neg$  (negation),  $\rightarrow$  (implication),  $\vee$  (disjunction),  $\wedge$  (conjunction), and  $\equiv$  (equivalence):

The connectives:  $\neg$  (negation),  $\rightarrow$  (implication),  $\vee$  (disjunction),  $\wedge$  (conjunction), and  $\equiv$  (equivalence):

|          |              |               |   |   |   |        |   |   |   |
|----------|--------------|---------------|---|---|---|--------|---|---|---|
| $\alpha$ | $\neg\alpha$ | $\rightarrow$ | f | u | t | $\vee$ | f | u | t |
| f        | t            | f             | t | t | t | f      | f | u | t |
| u        | u            | u             | u | u | t | u      | u | u | t |
| t        | f            | t             | f | u | t | t      | t | t | t |

|          |   |   |   |          |   |   |   |
|----------|---|---|---|----------|---|---|---|
| $\wedge$ | f | u | t | $\equiv$ | f | u | t |
| f        | f | f | f | f        | t | u | f |
| u        | f | u | u | u        | u | u | u |
| t        | f | u | t | t        | f | u | t |

The connectives:  $\neg$  (negation),  $\rightarrow$  (implication),  $\vee$  (disjunction),  $\wedge$  (conjunction), and  $\equiv$  (equivalence):

|          |              |
|----------|--------------|
| $\alpha$ | $\neg\alpha$ |
| f        | t            |
| u        | u            |
| t        | f            |

|               |   |   |   |
|---------------|---|---|---|
| $\rightarrow$ | f | u | t |
| f             | t | t | t |
| u             | u | u | t |
| t             | f | u | t |

|        |   |   |   |
|--------|---|---|---|
| $\vee$ | f | u | t |
| f      | f | u | t |
| u      | u | u | t |
| t      | t | t | t |

|          |   |   |   |
|----------|---|---|---|
| $\wedge$ | f | u | t |
| f        | f | f | f |
| u        | f | u | u |
| t        | f | u | t |

|          |   |   |   |
|----------|---|---|---|
| $\equiv$ | f | u | t |
| f        | t | u | f |
| u        | u | u | u |
| t        | f | u | t |

- $K_3 = (\{f, u, t\}, \neg, \rightarrow, \vee, \wedge, \{t\})$  *Kleene matrix*

- $K_3$  defines the non-tautological logic since any valuation which assigns the value  $u$  to each propositional variable sends any formula into  $u$

- $K_3$  defines the non-tautological logic since any valuation which assigns the value  $u$  to each propositional variable sends any formula into  $u$
- Kleene is aware of discarding even such law of logic as the law of identity  $p \rightarrow p$  and its “equivalential” version  $p \equiv p$ .

- $K_3$  defines the non-tautological logic since any valuation which assigns the value  $u$  to each propositional variable sends any formula into  $u$
- Kleene is aware of discarding even such law of logic as the law of identity  $p \rightarrow p$  and its “equivalential” version  $p \equiv p$ .
- ”Here “unknown” is a category into which we can regard any proposition as falling, whose value we either do not know or choose for the moment to disregard; and it does not then exclude the other two possibilities ‘true’ and ‘false’.” (Kleene [1952], p. 335)

- $K_3$  defines the non-tautological logic since any valuation which assigns the value  $u$  to each propositional variable sends any formula into  $u$
- Kleene is aware of discarding even such law of logic as the law of identity  $p \rightarrow p$  and its “equivalential” version  $p \equiv p$ .
- ”Here “unknown” is a category into which we can regard any proposition as falling, whose value we either do not know or choose for the moment to disregard; and it does not then exclude the other two possibilities ‘true’ and ‘false’.” (Kleene [1952], p. 335)
- We may, therefore, conclude that Kleene treated the added logical value as apparent or as pseudo-value, distinct from the real truth-values, cf. Turquette [1963].

## 2. Inexact classes

## 2. Inexact classes

Körner [1966] gave the most accurate and compatible interpretation of Kleene's views of the connectives of  $K3$ :

## 2. Inexact classes

Körner [1966] gave the most accurate and compatible interpretation of Kleene's views of the connectives of  $K3$ :

- He introduced a non-standard notion of an inexact class as an object characterized by a non-definite classifying procedure determined by a partial algorithm

## 2. Inexact classes

Körner [1966] gave the most accurate and compatible interpretation of Kleene's views of the connectives of  $K3$ :

- He introduced a non-standard notion of an inexact class as an object characterized by a non-definite classifying procedure determined by a partial algorithm
- An inexact class is an abstract class set by a "characteristic function", which discerns between members, non-members, and neutral candidates for membership

## 2. Inexact classes

Körner [1966] gave the most accurate and compatible interpretation of Kleene's views of the connectives of  $K3$ :

- He introduced a non-standard notion of an inexact class as an object characterized by a non-definite classifying procedure determined by a partial algorithm
- An inexact class is an abstract class set by a "characteristic function", which discerns between members, non-members, and neutral candidates for membership
- the algebra of inexact classes leads naturally to the three-valued logic  $K3$

## 2. Inexact classes

Körner [1966] gave the most accurate and compatible interpretation of Kleene's views of the connectives of  $K3$ :

- He introduced a non-standard notion of an inexact class as an object characterized by a non-definite classifying procedure determined by a partial algorithm
- An inexact class is an abstract class set by a "characteristic function", which discerns between members, non-members, and neutral candidates for membership
- the algebra of inexact classes leads naturally to the three-valued logic  $K3$
- Basing on it Cleave [1974] worked out a model-theoretic framework for the logic of inexact predicates and he formulated a notion of logical consequence.

Cleave's contribution

## Cleave's contribution

- is important in many respects, especially because of the analysis concerning the first-order inexact structures

## Cleave's contribution

- is important in many respects, especially because of the analysis concerning the first-order inexact structures
- is faulty since Cleave adopts the idea that the Kleene values  $f$ ,  $u$ ,  $t$  are mutually comparable and linearly ordered:  $f \leq u \leq t$ . (-1, 0 and +1 stand for  $f, u$  and  $t$ )

## Cleave's contribution

- is important in many respects, especially because of the analysis concerning the first-order inexact structures
- is faulty since Cleave adopts the idea that the Kleene values  $f$ ,  $u$ ,  $t$  are mutually comparable and linearly ordered:  $f \leq u \leq t$ . (-1, 0 and +1 stand for  $f, u$  and  $t$ )
- The concept of a a degree of truth preserving consequence operation is determined by the Cleave's variant of  $K3$ ,

## Cleave's contribution

- is important in many respects, especially because of the analysis concerning the first-order inexact structures
- is faulty since Cleave adopts the idea that the Kleene values  $f, u, t$  are mutually comparable and linearly ordered:  $f \leq u \leq t$ . (-1, 0 and +1 stand for  $f, u$  and  $t$ )
- The concept of a degree of truth preserving consequence operation is determined by the Cleave's variant of  $K3$ ,

$$C3 = ( -1, 0, +1 , \neg , \rightarrow , \vee , \wedge , \equiv , +1 ) ,$$

## Cleave's contribution

- is important in many respects, especially because of the analysis concerning the first-order inexact structures
- is faulty since Cleave adopts the idea that the Kleene values  $f, u, t$  are mutually comparable and linearly ordered:  $f \leq u \leq t$ . ( $-1, 0$  and  $+1$  stand for  $f, u$  and  $t$ )
- The concept of a a degree of truth preserving consequence operation is determined by the Cleave's variant of  $K3$ ,

$$C3 = ( -1, 0, +1 , \neg , \rightarrow , \vee , \wedge , \equiv , +1 ) ,$$

as follows:  $\alpha$  is a  $C3$  consequence of the set of formulas  $X$ ,  $X \models_{C3} \alpha$ , whenever for every valuation  $v$  of the language in  $K3$ ,  $\min\{v(\beta) : \beta \in X\} \leq v(\alpha)$ .

## Cleave's contribution

- is important in many respects, especially because of the analysis concerning the first-order inexact structures
- is faulty since Cleave adopts the idea that the Kleene values  $f, u, t$  are mutually comparable and linearly ordered:  $f \leq u \leq t$ . ( $-1, 0$  and  $+1$  stand for  $f, u$  and  $t$ )
- The concept of a degree of truth preserving consequence operation is determined by the Cleave's variant of  $K3$ ,

$$C3 = ( -1, 0, +1, \neg, \rightarrow, \vee, \wedge, \equiv, +1 ),$$

as follows:  $\alpha$  is a  $C3$  consequence of the set of formulas  $X$ ,  $X \models_{C3} \alpha$ , whenever for every valuation  $v$  of the language in  $K3$ ,  $\min\{v(\beta) : \beta \in X\} \leq v(\alpha)$ .

### 3. Inference

### **3. Inference**

Departing from the idea of "neutrality" of sentences of the third category and accepting inexactness, Körner [1966] tries to apply the classical logic to empirical discourse:

### 3. Inference

Departing from the idea of "neutrality" of sentences of the third category and accepting inexactness, Körner [1966] tries to apply the classical logic to empirical discourse:

- To this aim, he proposes a special evaluation device leading to a kind of "modified two-valued logic" and, ultimately, to an instrument of evaluation validity of sentences and deduction

### 3. Inference

Departing from the idea of "neutrality" of sentences of the third category and accepting inexactness, Körner [1966] tries to apply the classical logic to empirical discourse:

- To this aim, he proposes a special evaluation device leading to a kind of "modified two-valued logic" and, ultimately, to an instrument of evaluation validity of sentences and deduction
- His interesting trial ends negatively: "The classical two-valued logic as an instrument of deduction, however, presupposes that neutral propositions are treated as if they were true, and inexact predicates as if they were exact. ..."

### 3. Inference

Departing from the idea of "neutrality" of sentences of the third category and accepting inexactness, Körner [1966] tries to apply the classical logic to empirical discourse:

- To this aim, he proposes a special evaluation device leading to a kind of "modified two-valued logic" and, ultimately, to an instrument of evaluation validity of sentences and deduction
- His interesting trial ends negatively: "The classical two-valued logic as an instrument of deduction, however, presupposes that neutral propositions are treated as if they were true, and inexact predicates as if they were exact. ..."
- This obviously means that the use of the classical logic for drawing empirical conclusions from empirical premises is unsound.

## Grounds for new Kleene inference relation

- Körner's analysis of the structure of scientific theories and their relation to the experience, including:

## Grounds for new Kleene inference relation

- Körner's analysis of the structure of scientific theories and their relation to the experience, including:
- K1. Two complementary spheres of scientific activity: the sphere of experience and the sphere of theory

## Grounds for new Kleene inference relation

- Körner's analysis of the structure of scientific theories and their relation to the experience, including:
- K1. Two complementary spheres of scientific activity: the sphere of experience and the sphere of theory
- K2. Deductive unification of a field of experience by classical logic is an idealisation or a modification of empirical discourse

## Grounds for new Kleene inference relation

- Körner's analysis of the structure of scientific theories and their relation to the experience, including:
- K1. Two complementary spheres of scientific activity: the sphere of experience and the sphere of theory
- K2. Deductive unification of a field of experience by classical logic is an idealisation or a modification of empirical discourse
- K3. Idealization procedure supported by the two operators which convert neutral propositions into false or true

## Grounds for new Kleene inference relation

- Körner's analysis of the structure of scientific theories and their relation to the experience, including:
- K1. Two complementary spheres of scientific activity: the sphere of experience and the sphere of theory
- K2. Deductive unification of a field of experience by classical logic is an idealisation or a modification of empirical discourse
- K3. Idealization procedure supported by the two operators which convert neutral propositions into false or true

The relation of empirical inference:

- holds between true or indefinite premises and the true conclusion. Thus, a conclusion  $\alpha$  may be empirically inferred from a set of premises  $X$ , whenever, for any interpretation, it is the case that if all elements of  $X$  are not false, then  $\alpha$  is true

## The relation of empirical inference:

- holds between true or indefinite premises and the true conclusion. Thus, a conclusion  $\alpha$  may be empirically inferred from a set of premises  $X$ , whenever, for any interpretation, it is the case that if all elements of  $X$  are not false, then  $\alpha$  is true
- may express the passing from the empirical inexact and indefinite discourse to the "heaven" of theory but at the cost of a natural extension of the language

## The relation of empirical inference:

- holds between true or indefinite premises and the true conclusion. Thus, a conclusion  $\alpha$  may be empirically inferred from a set of premises  $X$ , whenever, for any interpretation, it is the case that if all elements of  $X$  are not false, then  $\alpha$  is true
- may express the passing from the empirical inexact and indefinite discourse to the "heaven" of theory but at the cost of a natural extension of the language
- the latter is possible when two connectives corresponding to idealisation procedures are used. In case of (inexact) classes this means replacement by their exact counterparts, in case of sentences turning indefinite sentences into definite, i.e. true or false

## The relation of empirical inference:

- holds between true or indefinite premises and the true conclusion. Thus, a conclusion  $\alpha$  may be empirically inferred from a set of premises  $X$ , whenever, for any interpretation, it is the case that if all elements of  $X$  are not false, then  $\alpha$  is true
- may express the passing from the empirical inexact and indefinite discourse to the "heaven" of theory but at the cost of a natural extension of the language
- the latter is possible when two connectives corresponding to idealisation procedures are used. In case of (inexact) classes this means replacement by their exact counterparts, in case of sentences turning indefinite sentences into definite, i.e. true or false

Definition:

Definition:

A relation  $\vdash_{K3}$  is said to be a matrix empirical inference of K3 provided that for any  $X \subseteq For, \alpha \in For$

Definition:

A relation  $\vdash_{K3}$  is said to be a matrix empirical inference of K3 provided that for any  $X \subseteq For, \alpha \in For$

$X \vdash_{K3} \alpha$  iff for every  $h \in Hom(L, A)(h(\alpha) = t$  whenever  $h(X) \subseteq u, t$ .

Definition:

A relation  $\vdash_{K3}$  is said to be a matrix empirical inference of K3 provided that for any  $X \subseteq For, \alpha \in For$

$X \vdash_{K3} \alpha$  iff for every  $h \in Hom(L, A)(h(\alpha) = t$  whenever  $h(X) \subseteq u, t$ .

- An intuition behind this definition is that a conclusion may be inferred empirically from a set of premises  $X$ , when for any interpretation it is the case that if all elements of  $X$  are not false then  $\alpha$  is true

Definition:

A relation  $\vdash_{K3}$  is said to be a matrix empirical inference of K3 provided that for any  $X \subseteq For, \alpha \in For$

$X \vdash_{K3} \alpha$  iff for every  $h \in Hom(L, A)(h(\alpha) = t$  whenever  $h(X) \subseteq u, t$ .

- An intuition behind this definition is that a conclusion may be inferred empirically from a set of premises  $X$ , when for any interpretation it is the case that if all elements of  $X$  are not false then  $\alpha$  is true
- $\vdash_{K3}$  is compatible with two possible logical idealisations of undefinite empirical sentences

Definition:

A relation  $\vdash_{K3}$  is said to be a matrix empirical inference of K3 provided that for any  $X \subseteq For, \alpha \in For$

$X \vdash_{K3} \alpha$  iff for every  $h \in Hom(L, A)(h(\alpha) = t$  whenever  $h(X) \subseteq u, t$ .

- An intuition behind this definition is that a conclusion may be inferred empirically from a set of premises  $X$ , when for any interpretation it is the case that if all elements of  $X$  are not false then  $\alpha$  is true
- $\vdash_{K3}$  is compatible with two possible logical idealisations of undefinite empirical sentences

The operators corresponding to these idealisations are presented through Kleene-like tables as follows:

The operators corresponding to these idealisations are presented through Kleene-like tables as follows:

|          |           |          |           |
|----------|-----------|----------|-----------|
| $\alpha$ | $W\alpha$ | $\alpha$ | $S\alpha$ |
| f        | f         | f        | f         |
| u        | f         | u        | t         |
| t        | t         | t        | t         |

The operators corresponding to these idealisations are presented through Kleene-like tables as follows:

| $\alpha$ | $W\alpha$ |
|----------|-----------|
| f        | f         |
| u        | f         |
| t        | t         |

| $\alpha$ | $S\alpha$ |
|----------|-----------|
| f        | f         |
| u        | t         |
| t        | t         |

The letters chosen for these connectives may be read as Weak and Strong. Let us note that the empirical inference just defined may be also interpreted as a relation which holds between a set of premises  $X$  and a conclusion whenever it is true independently from accepted idealisation(s) of empirical sentences in  $X$ .

Here are examples of theorems establishing some essential properties of the empirical inference of Kleene sentential logic:

Here are examples of theorems establishing some essential properties of the empirical inference of Kleene sentential logic:

- $(e_1)$  If  $X \vdash_{K3} \alpha$  and  $X \vdash_{K3} \alpha \rightarrow \beta$ , then  $X \vdash_{K3} \beta$

Here are examples of theorems establishing some essential properties of the empirical inference of Kleene sentential logic:

- $(e_1)$  If  $X \vdash_{K3} \alpha$  and  $X \vdash_{K3} \alpha \rightarrow \beta$ , then  $X \vdash_{K3} \beta$
- $(e_2)$   $X \vdash_{K3} \alpha \rightarrow \beta$  if and only if  $X \cup \alpha \vdash_{K3} \beta$

Here are examples of theorems establishing some essential properties of the empirical inference of Kleene sentential logic:

- $(e_1)$  If  $X \vdash_{K3} \alpha$  and  $X \vdash_{K3} \alpha \rightarrow \beta$ , then  $X \vdash_{K3} \beta$
- $(e_2)$   $X \vdash_{K3} \alpha \rightarrow \beta$  if and only if  $X \cup \alpha \vdash_{K3} \beta$
- $(e_3)$   $\alpha \vdash_{K3} S(\alpha)$  and  $W(\alpha) \vdash_{K3} \alpha$ .

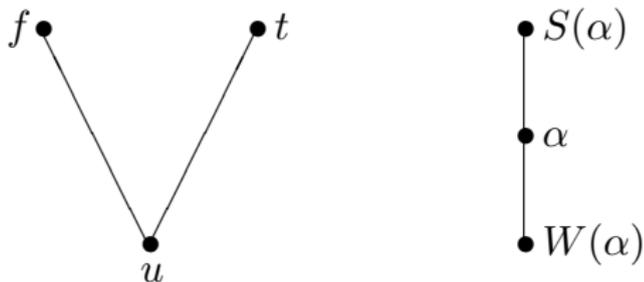
Here are examples of theorems establishing some essential properties of the empirical inference of Kleene sentential logic:

- $(e_1)$  If  $X \vdash_{K3} \alpha$  and  $X \vdash_{K3} \alpha \rightarrow \beta$ , then  $X \vdash_{K3} \beta$
- $(e_2)$   $X \vdash_{K3} \alpha \rightarrow \beta$  if and only if  $X \cup \alpha \vdash_{K3} \beta$
- $(e_3)$   $\alpha \vdash_{K3} S(\alpha)$  and  $W(\alpha) \vdash_{K3} \alpha$ .

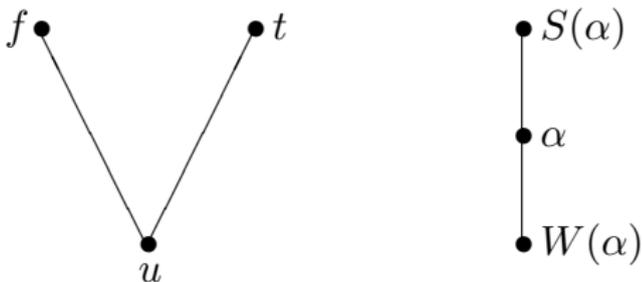
The first property is an *inferential modus ponens*, the second a *conditional deduction theorem*. The inferences in  $(e_3)$  state that from a sentence its strong idealisation follows while the sentence itself is inferred from its weak idealisation.

The diagram below shows how the Kleene values  $f, u, t$  are ordered and displays the inferential relations between a formula  $\alpha$  and its two idealizations,  $S(\alpha)$  and  $W(\alpha)$ :

The diagram below shows how the Kleene values  $f, u, t$  are ordered and displays the inferential relations between a formula  $\alpha$  and its two idealizations,  $S(\alpha)$  and  $W(\alpha)$ :



The diagram below shows how the Kleene values  $f, u, t$  are ordered and displays the inferential relations between a formula  $\alpha$  and its two idealizations,  $S(\alpha)$  and  $W(\alpha)$ :



So, any sentence  $\alpha$  is deductively situated between its two idealizations:  $W(\alpha) \vdash_{K3} \alpha \vdash_{K3} S(\alpha)$ .

## 4. Final remarks

#### 4. Final remarks

- The empirical inference relation appears to be a special case of the semantic relation defined on q-matrices  $M^* = (A, D^*, D)$ , where  $D^*, D$  are disjoint subsets of  $A$  interpreted as rejected and accepted elements.

## 4. Final remarks

- The empirical inference relation appears to be a special case of the semantic relation defined on q-matrices  $M^* = (A, D^*, D)$ , where  $D^*, D$  are disjoint subsets of  $A$  interpreted as rejected and accepted elements.
- Kleene's case:  $A = (\{f, u, t\}, \neg, \rightarrow, \vee, \wedge, \equiv, S, W)$ ,  $D^* = \{f\}$  and  $D = \{t\}$ .

#### 4. Final remarks

- The empirical inference relation appears to be a special case of the semantic relation defined on q-matrices  $M^* = (A, D^*, D)$ , where  $D^*, D$  are disjoint subsets of  $A$  interpreted as rejected and accepted elements.
- Kleene's case:  $A = (\{f, u, t\}, \neg, \rightarrow, \vee, \wedge, \equiv, S, W)$ ,  $D^* = \{f\}$  and  $D = \{t\}$ .
- The Kleene q-matrix consequence is a  $W$  and  $S$  compatible inference from non-rejected premises to accepted conclusions.

## 2. References

- [1] Cleave, J.P., The notion of logical consequence in the logic of inexact predicates, **Zeitschrift für Mathematische Logik and Grundlagen der Mathematik**, 20, 1974, 307-324.
- [2] Kleene, S.C., 'On a notation for ordinal numbers', **The Journal of Symbolic Logic**, 3(1938), pp. 150–155.
- [3] Kleene, S.C., **Introduction to metamathematics**, North-Holland, Amsterdam, 1952.
- [4] Körner, S., **Experience and theory**, Routledge and Kegan Paul, London, 1966.
- [5] Malinowski, G., Inferential many-valuedness [in:] Woleński, J. (ed.) **Philosophical logic in Poland**, Synthese Library, 228, Kluwer Academic Publishers, Dordrecht, 1994, 75-84.
- [6] Malinowski, G., **Many-valued logics**, Oxford Logic Guides, 25, Clarendon Press, Oxford, 1993.

- [7] Malinowski, G., Q-consequence operation, **Reports on Mathematical Logic**, 24, 1990, 49-59.
- [8] Turquette, A. R., Modality, minimality, and many-valuedness, **Acta Philosophica Fennica**, XVI, 1963, 261-276.

Department of Logic  
University of Łódź  
Poland  
*gregmal@uni.lodz.pl*